

DP-SPRT: Differentially Private Sequential Hypothesis Testing

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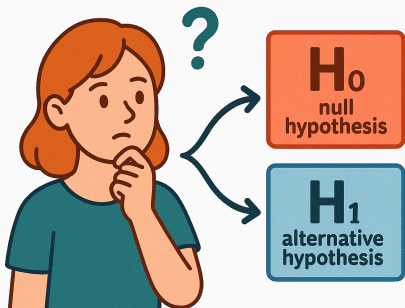


Sequential Hypothesis Testing

Setup: We observe samples $X_1, X_2, \dots$ sequentially from distribution ν_θ

Goal: Test $H_0 : \theta = \theta_0$ vs $H_1 : \theta = \theta_1$ with as few samples as possible

Constraints: False positive probability $\leq \alpha$, false negative probability $\leq \beta$

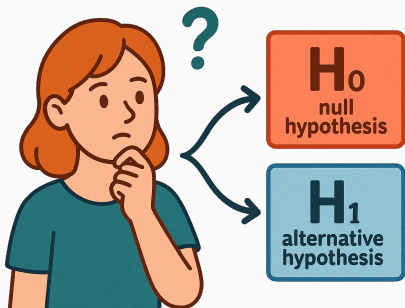


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When do we need to test ?

- Clinical Trials
- A/B Testing
- Fraud Detection

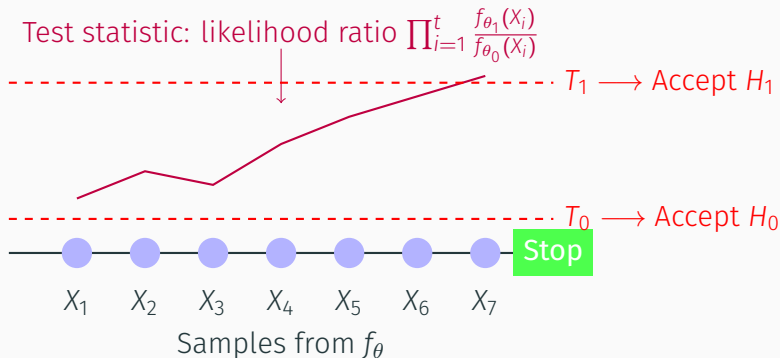
Sequential Probability Ratio Test

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Sequential Probability Ratio Test (SPRT) (Wald, 1945):



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SPRT for one dimensionnal exponential families: Stop at $\tau = \min(\tau_0, \tau_1)$ with decision $\hat{d} = i$ if $\tau = \tau_i$ where

$$\tau_0 = \inf \left\{ n : \bar{X}_n \leq \mu_0 + \frac{\text{KL}(\nu_{\theta_0}, \nu_{\theta_1}) - \log(1/\beta)/n}{\theta_1 - \theta_0} \right\}$$
$$\tau_1 = \inf \left\{ n : \bar{X}_n \geq \mu_1 - \frac{\text{KL}(\nu_{\theta_1}, \nu_{\theta_0}) - \log(1/\alpha)/n}{\theta_1 - \theta_0} \right\}.$$

$\text{KL}(\nu_{\theta_0}, \nu_{\theta_1})$ is the KL divergence between the two distributions and $\bar{X}_n$ is the empirical mean of the samples up to time n .

The Privacy Problem: A Medical Trial

Scenario: Testing if new drug works better than placebo

H_0 : Drug success rate = 30% (Same as placebo) vs H_1 : Drug success rate = 70%

Each patient outcome: $X_i \in \{0, 1\}$ (failure/success)



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Question: What was patient 7's outcome? **Success (1) or Failure (0)?**

The stopping decision reveals patient 7 had a SUCCESS!

Privacy violation: Patient 7's medical outcome is leaked by our decision to stop

Privacy in Sequential Decisions



Clinical Trial

- Testing new drug vs placebo
- **Stopping pattern** reveals:
 - Treatment effectiveness
 - Patient responses



A/B Testing

- Users see different versions
- **When we conclude** reveals:
 - User behavior
 - Conversion rates



Fraud Detection

- Monitor transactions
- **Alert timing** reveals:
 - Transaction patterns
 - Detection methods

Core Problem: When we stop reveals what we observed

Differential Privacy

Neighboring Datasets: Two datasets D, D' are neighboring if they differ by exactly one record.

Differential Privacy (Dwork, Roth, et al., 2014): A randomized mechanism $\mathcal{M}$ is DP if for any neighboring datasets D and D' and for any events S :

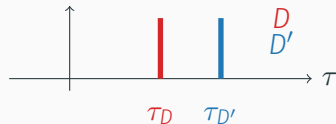
$$\epsilon\text{-DP: } \log \left(\frac{\mathbb{P}[\mathcal{M}(D) \in S]}{\mathbb{P}[\mathcal{M}(D') \in S]} \right) \leq \epsilon$$

$$(\alpha, \epsilon)\text{-Rényi DP: } D_\alpha(\mathcal{M}(D) \parallel \mathcal{M}(D')) \leq \epsilon$$

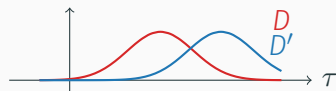
where D_α is the Rényi divergence of order $\alpha > 1$.

Stopping Time Distributions

Deterministic Algorithm



Private Algorithm



Privacy adds randomness to mask the stopping decision pattern

Our Method: DP-SPRT

DP-SPRT Algorithm (blue = privacy additions to SPRT):

Input: Hypotheses θ_0, θ_1 , error probabilities α, β , noise distributions $\mathcal{D}_Z, \mathcal{D}_Y$, correction function $C(n, x)$, error allocation γ

1. Sample threshold noise $Z \sim \mathcal{D}_Z$
2. **For** $n = 1, 2, 3, \dots$ **do**
3. Sample query noise $Y_n \sim \mathcal{D}_Y$
4. Compute noisy average $\bar{X}_n = \frac{1}{n} \sum_{i=1}^n X_i + \frac{Y_n}{n}$
5. Compute noisy threshold $\hat{T}_0^n = \mu_0 + \frac{\text{KL}(\nu_{\theta_0}, \nu_{\theta_1}) - \log(1/(\gamma\beta))/n}{\theta_1 - \theta_0} - C(n, (1 - \gamma)\beta) - \frac{Z}{n}$
6. Compute noisy threshold $\hat{T}_1^n = \mu_1 - \frac{\text{KL}(\nu_{\theta_1}, \nu_{\theta_0}) - \log(1/(\gamma\alpha))/n}{\theta_1 - \theta_0} + C(n, (1 - \gamma)\alpha) + \frac{Z}{n}$
7. **If** noisy average $\bar{X}_n$ is below noisy threshold $\hat{T}_0^n$ **then** Halt and accept $\mathcal{H}_0$
8. **Else if** noisy average $\bar{X}_n$ is above noisy threshold $\hat{T}_1^n$ **then** Halt and accept $\mathcal{H}_1$

Y_n and Z are used for both conditions unlike when composing AboveThreshold

Table: Comparison of DP-SPRT instantiations

	Laplace Noise	Gaussian Noise
Noise distributions	$Y_n \sim \text{Lap}(4/\varepsilon)$ $Z \sim \text{Lap}(2/\varepsilon)$	$Y_n \sim \mathcal{N}(0, \sigma_Y^2)$ $Z \sim \mathcal{N}(0, \sigma_Z^2)$
Privacy guarantee	ε -Differential Privacy	$(\alpha, \varepsilon(\alpha))$ -RDP
Correction function $C(n, x)$	$\frac{6 \log(n^s \zeta(s)/x)}{n\varepsilon}$	$\frac{\sqrt{2(\sigma_Y^2 + \sigma_Z^2) \log(n^s \zeta(s)/2)}}{n}$

Correctness: DP-SPRT satisfies $\mathbb{P}_{\theta_0}(\hat{d} = 1) \leq \alpha$ and $\mathbb{P}_{\theta_1}(\hat{d} = 0) \leq \beta$ when the distribution $\mathcal{D}_Y$ is symmetric and the correction function $C(n, x)$ satisfies:

$$\sum_{n=1}^{\infty} \mathbb{P} \left(\frac{Y_n}{n} - \frac{Z}{n} > C(n, x) \right) \leq x$$

Near-Optimal Sample Complexity (DP-SPRT with Laplace noise)

Lower Bound (any ε -DP test):

$$\mathbb{E}[\tau] \geq \frac{\log(1/\beta)}{\min(\text{KL}(\nu_{\theta_0}, \nu_{\theta_1}), \varepsilon \cdot \text{TV}(\nu_{\theta_0}, \nu_{\theta_1}))}$$

where TV is the total variation distance between two distributions.

Sample Complexity Upper Bound (Laplace Noise):

$$\mathbb{E}[\tau] \lesssim \max \left(\frac{\log(1/\beta)}{\text{KL}(\nu_{\theta_0}, \nu_{\theta_1})}, \frac{(\theta_1 - \theta_0) \log(1/\beta)}{\varepsilon \cdot \text{KL}(\nu_{\theta_0}, \nu_{\theta_1})} \right)$$

For Bernoulli distributions, we have

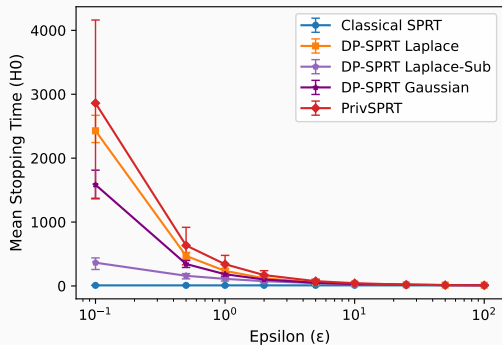
$$\frac{\text{KL}(\nu_{\theta_0}, \nu_{\theta_1})}{\theta_1 - \theta_0} \xrightarrow{\theta_1 \rightarrow \theta_0} \text{TV}(\nu_{\theta_0}, \nu_{\theta_1})$$

DP-SPRT with Laplace noise is near-optimal

Experimental Results: Performance Comparison

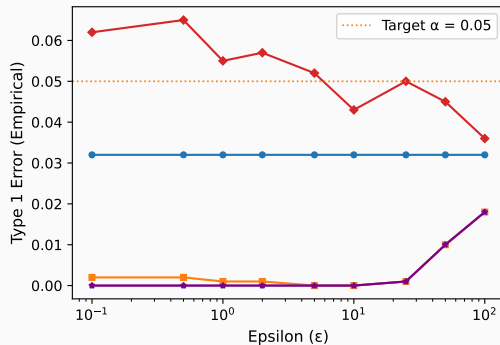
Setup: Bernoulli($p_0 = 0.3$) vs Bernoulli($p_1 = 0.7$), $\alpha = \beta = 0.05$, 1000 trials

Sample Complexity vs Privacy Level



On average, **DP-SPRT variants outperform** PrivSPRT (Zhang, Mei, and Cummings, 2022) across privacy levels

Error Control Comparison



All DP-SPRT variants guarantee error control, while PrivSPRT can violate error targets due to empirical tuning

Conclusion

Privacy:

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Our Contributions:

1. **Theoretically calibrated** private sequential test with **guaranteed error control**
2. **Near-optimal sample complexity** matching lower bounds up to a constant in some regimes
3. **Practical implementation** with **no empirical tuning** required, **low variance** in stopping times, and **subsampling amplification** in high-privacy regimes

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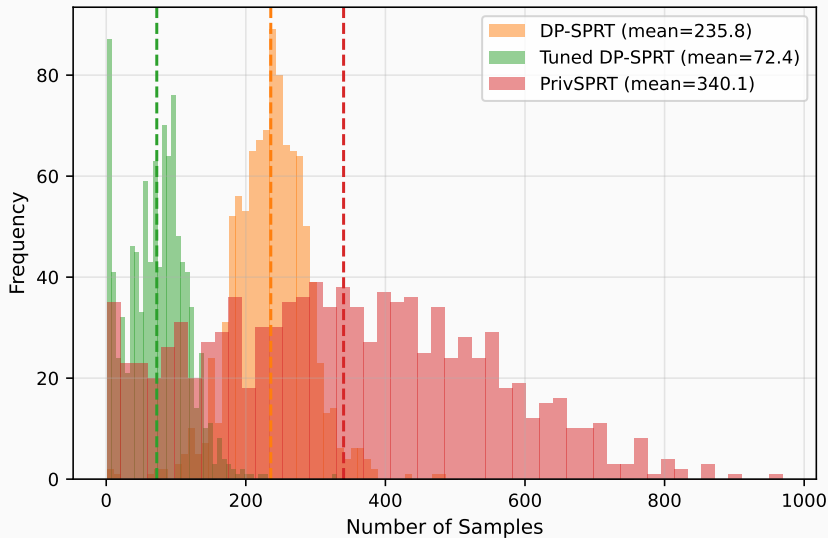
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Thank you! Questions?



-  Dwork, Cynthia, Aaron Roth, et al. (2014). **“The algorithmic foundations of differential privacy”**. In: *Foundations and Trends® in Theoretical Computer Science* 9.3–4, pp. 211–407.
-  Wald, A. (1945). **“Sequential Tests of Statistical Hypotheses”**. In: *Annals of Mathematical Statistics* 16(2), pp. 117–186.
-  Zhang, Wanrong, Yajun Mei, and Rachel Cummings (2022). **“Private Sequential Hypothesis Testing for Statisticians: Privacy, Error Rates, and Sample Size”**. In: *AISTATS*.

Distribution of Stopping Time



Theorem (Privacy)

Let $\mathcal{A}$ be the DP-SPRT algorithm and τ the (random) stopping time.

(i) ϵ -DP: If noise mechanisms $\mathcal{D}_Z$ and $\mathcal{D}_Y$ satisfy ϵ_Z -DP (sensitivity 1) and ϵ_Y -DP (sensitivity 2), then DP-SPRT satisfies $(\epsilon_Z + \epsilon_Y)$ -DP.

(ii) (α, ϵ) -RDP: If mechanisms have RDP profiles $\epsilon_Z(\alpha)$ and $\epsilon_Y(\alpha)$, then:

$$\mathbb{D}_\alpha(\mathcal{A}(D) \parallel \mathcal{A}(D')) \leq \frac{\alpha - 1/2}{\alpha - 1} \epsilon_Z(2\alpha) + \epsilon_Y(\alpha) + \frac{\log \left(2 \mathbb{E}_{Z \sim \mathcal{D}_Z} [\mathbb{E}_{(\tau, \hat{d}) \sim \mathcal{A}(D')} [\tau | Z|^2]] \right)}{2(\alpha - 1)}$$

Theorem (Correctness)

For any error allocation $\gamma \in (0, 1)$, DP-SPRT is (α, β) -correct if:

$$\forall \delta \in (0, 1), \quad \sum_{n=1}^{\infty} \mathbb{P} \left(\frac{Y_n}{n} - \frac{Z}{n} > C(n, \delta) \right) \leq \delta$$

Proof idea: Union bound decomposes error into SPRT error + privacy noise error

Theorem (Sample Complexity Upper Bound)

Assume correction condition holds. For error allocation $\gamma \in (0, 1)$ and $i \in \{0, 1\}$:

$$\mathbb{E}_{\theta_0}[\tau] \leq 1 + (1 - \gamma)\beta + \frac{1}{1 - \exp\left(-\frac{(TV(\nu_{\theta_0}, \nu_{\theta_1}))^4}{2(\theta_1 - \theta_0)^2}\right)} + N(\theta_0, \theta_1, \beta, \gamma)$$

$$\mathbb{E}_{\theta_1}[\tau] \leq 1 + (1 - \gamma)\alpha + \frac{1}{1 - \exp\left(-\frac{(TV(\nu_{\theta_0}, \nu_{\theta_1}))^4}{2(\theta_1 - \theta_0)^2}\right)} + N(\theta_1, \theta_0, \alpha, \gamma)$$

where $N(\theta, \theta', \delta, \gamma) = \inf \left\{ n : \frac{\log(1/(\delta\gamma))/n}{|\theta - \theta'|} + 2C(n, (1 - \gamma)\delta) \leq \frac{1}{2} \frac{KL(\nu_\theta, \nu_{\theta'})}{|\theta - \theta'|} \right\}$

Theorem (Lower Bound)

Any ε -DP test with $\mathbb{P}_{\theta_0}(\hat{d} = 1) \leq \alpha$, $\mathbb{P}_{\theta_1}(\hat{d} = 0) \leq \beta$ satisfies:

$$\mathbb{E}_{\theta_0}[\tau] \geq \frac{kl(\alpha, 1 - \beta)}{\min(KL(\nu_{\theta_0}, \nu_{\theta_1}), \varepsilon \cdot TV(\nu_{\theta_0}, \nu_{\theta_1}))}$$
$$\mathbb{E}_{\theta_1}[\tau] \geq \frac{kl(\beta, 1 - \alpha)}{\min(KL(\nu_{\theta_1}, \nu_{\theta_0}), \varepsilon \cdot TV(\nu_{\theta_0}, \nu_{\theta_1}))}$$

where $kl(x, y) = x \log(x/y) + (1 - x) \log((1 - x)/(1 - y))$ is binary relative entropy.

Corollary (Laplace DP-SPRT Sample Complexity)

For DP-SPRT with Laplace noise ($Y_n \sim \text{Lap}(4/\varepsilon)$, $Z \sim \text{Lap}(2/\varepsilon)$):

$$N(\theta_0, \theta_1, \beta, \gamma) \leq \frac{2 \log(1/(\gamma\beta))}{\text{KL}(\nu_{\theta_0}, \nu_{\theta_1})} + \frac{24(\theta_1 - \theta_0) \log(\zeta(s)/(1 - \gamma)\beta)}{\varepsilon \text{KL}(\nu_{\theta_0}, \nu_{\theta_1})} + o_{\beta \rightarrow 0}(\log(1/\beta))$$

Two Regimes:

- **Statistics-dominated:** ε large $\Rightarrow$ complexity $\approx \frac{\log(1/\beta)}{\text{KL}}$ (like non-private SPRT)
- **Privacy-dominated:** ε small $\Rightarrow$ complexity $\approx \frac{(\theta_1 - \theta_0) \log(1/\beta)}{\varepsilon \cdot \text{KL}}$